Rectangle-of-influence triangulations

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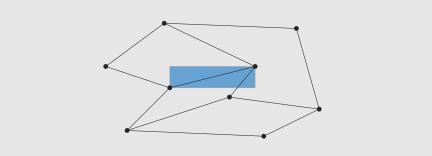
RI-Edges

• An edge is <u>RI</u> if its supporting rectangle (smallest axis-aligned bounding box) is empty of (other) points



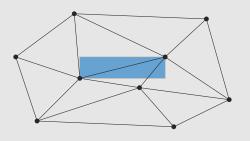
RI-Drawings

- · Drawing of a graph where all edges are RI
- Well-studied in Graph Drawing community



RI-Triangulations

- All internal faces are triangles
- Maximal



RI-Problems

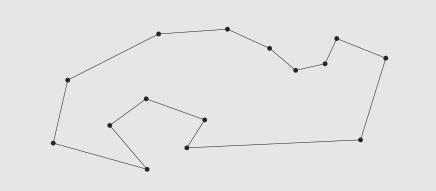
- 1. RI-triangulating a polygon
- 2. RI-triangulating a point set
- 3. Flipping one RI-triangulation to another
- 4. Flipping a triangulation to an RI-triangulation

RI-Problems

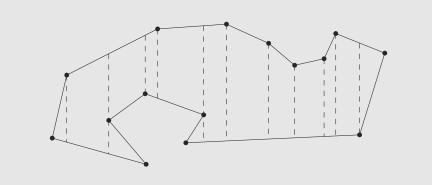
1. RI-triangulating a polygon

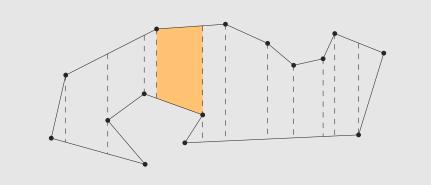
- 2. RI-triangulating a point set
- 3. Flipping one RI-triangulation to another
- 4. Flipping a triangulation to an RI-triangulation

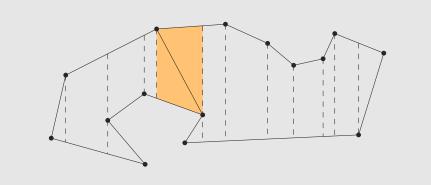
• All edges are RI

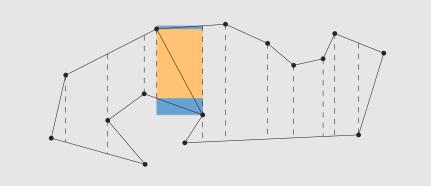


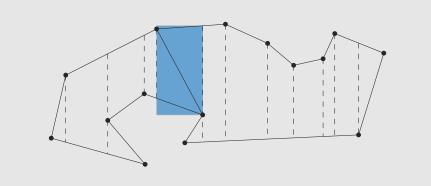
Compute trapezoidal decomposition

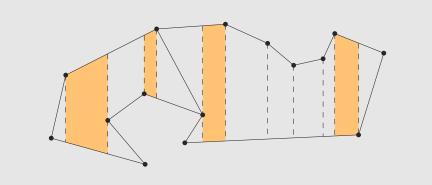


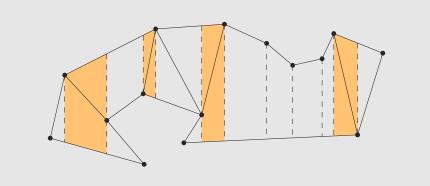


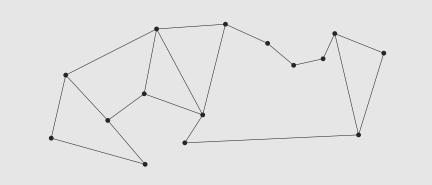




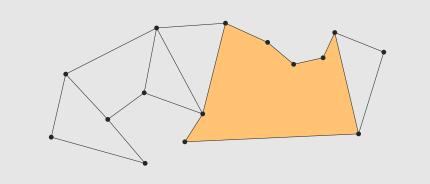


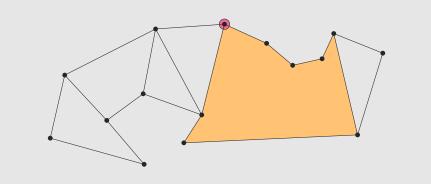


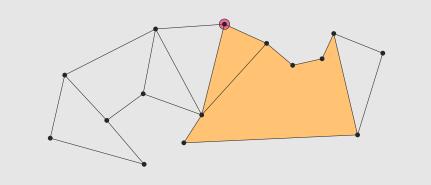


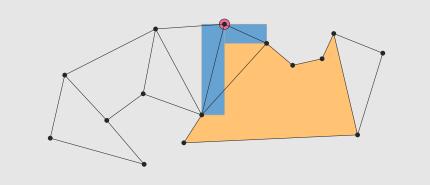


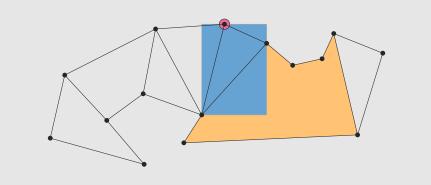
• Remaining pieces are x-monotone and one-sided

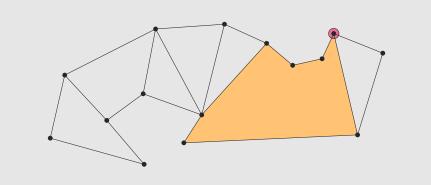


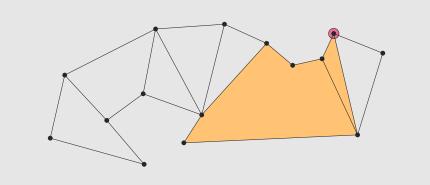


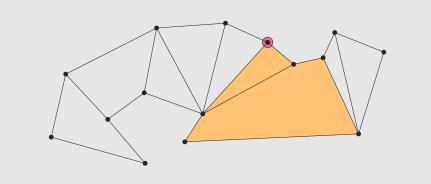


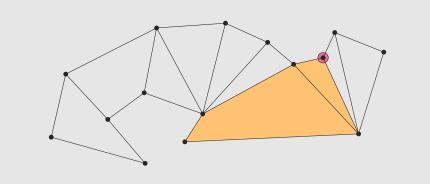


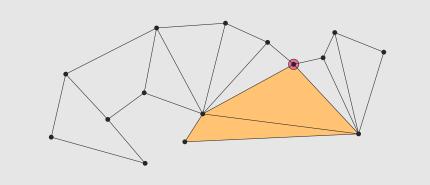


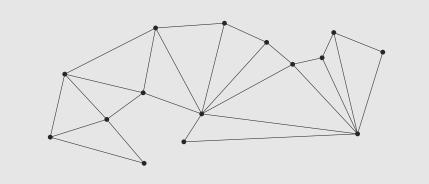






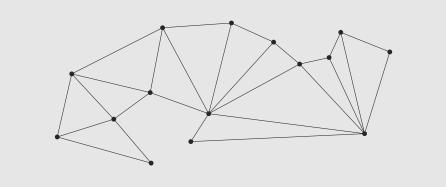






Theorem

Every RI-polygon can be RI-triangulated in linear time.

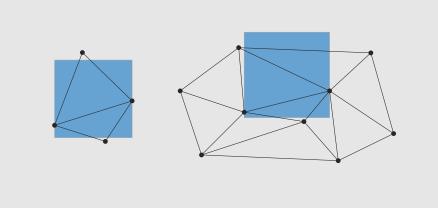


RI-Problems

- 1. RI-triangulating a polygon \checkmark
- 2. RI-triangulating a point set
- 3. Flipping one RI-triangulation to another
- 4. Flipping a triangulation to an RI-triangulation

RI-Point Sets

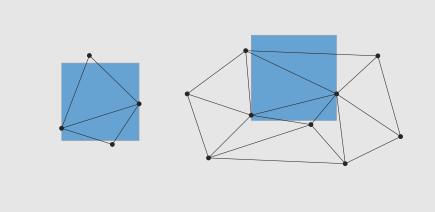
• The L_{∞} -Delaunay triangulation is RI



RI-Point Sets

Theorem

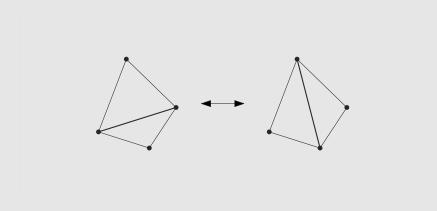
Any point set can be RI-triangulated in $O(n \log n)$ time.



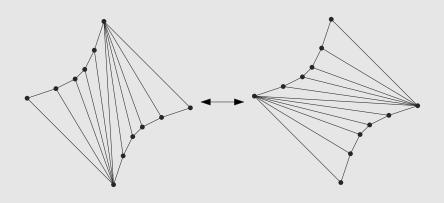
RI-Problems

- 1. RI-triangulating a polygon \checkmark
- 2. RI-triangulating a point set \checkmark
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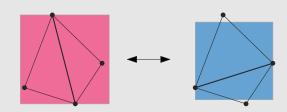
· Exchange one diagonal of a convex quad for the other



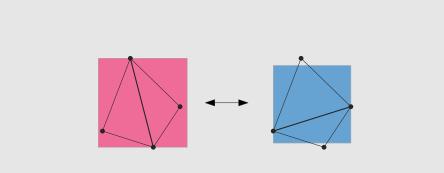
- · Is the class of RI-triangulations closed under flips?
- Diameter is $\Omega(n^2)$



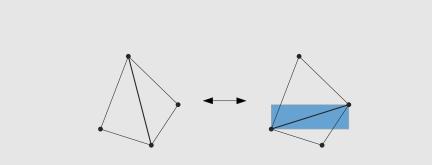
- Transform into the L_∞ -Delaunay triangulation
- An edge is locally L_∞ if its is L_∞ w.r.t. its neighbouring triangles
- If all edges are locally $L_\infty,$ we are in the $L_\infty\text{-}\textsc{Delaunay}$ triangulation



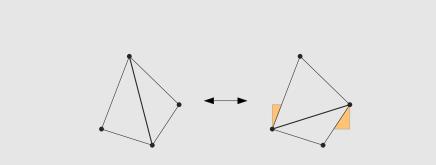
- Flip edges that are not locally L_∞
- · How do we know that new edge is (globally) RI?



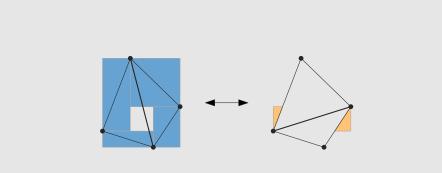
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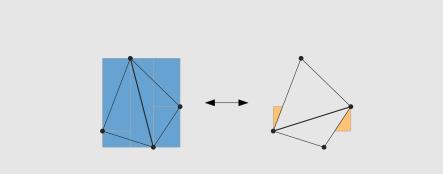
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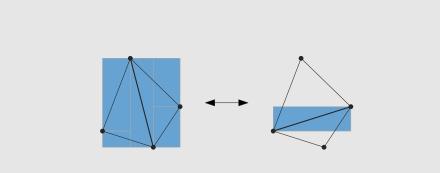
- Flip edges that are not locally L_∞
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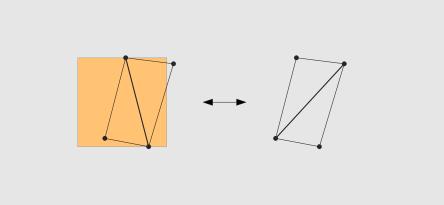
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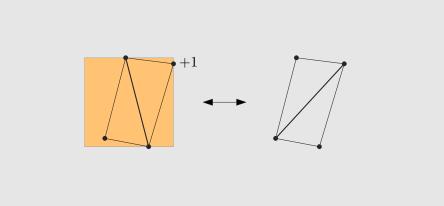
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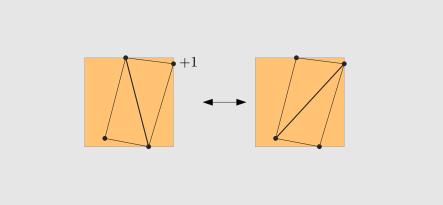
- · How many flips do we need?
- Give every edge a supporting square and count points inside



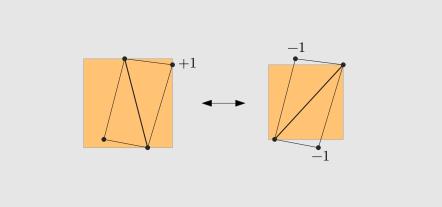
- · How many flips do we need?
- Give every edge a supporting square and count points inside



- · How many flips do we need?
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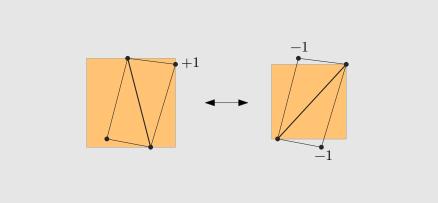


- · How many flips do we need?
- Give every edge a supporting square and count points inside



Theorem

The class of RI-triangulations is closed under flips and its diameter is $\Theta(n^2).$



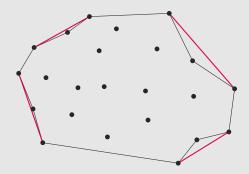
RI-Problems

- 1. RI-triangulating a polygon \checkmark
- 2. RI-triangulating a point set \checkmark
- 3. Flipping one RI-triangulation to another \checkmark

4. Flipping a triangulation to an RI-triangulation

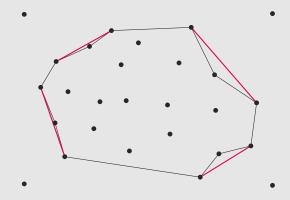
RI-Point Sets

The outer face can be messy



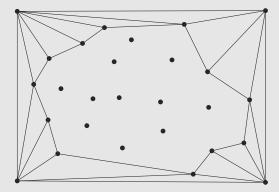
RI-Point Sets

- The outer face can be messy
- · We add 4 points 'far away' to deal with this

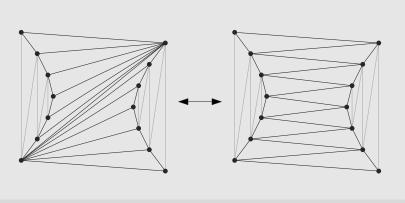


RI-Point Sets

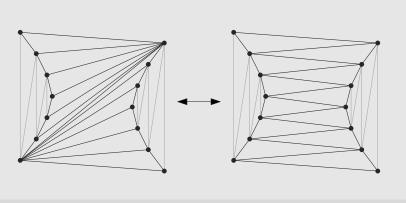
- The outer face can be messy
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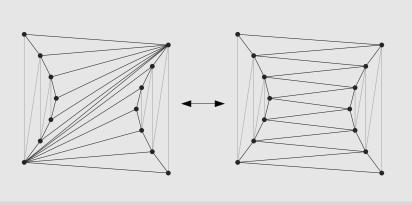
· Can we flip an arbitrary triangulation into an RI one?



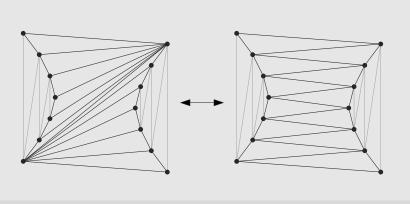
- · Can we flip an arbitrary triangulation into an RI one?
- Any triangulation can be flipped to any other in $O(n^2)$ flips [Lawson, 1972]



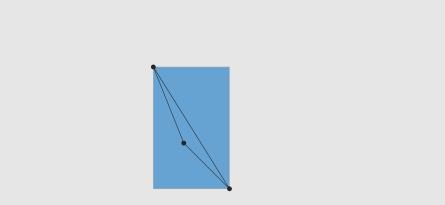
- · Can we flip an arbitrary triangulation into an RI one?
- Any triangulation can be flipped to any other in $O(n^2)$ flips [Lawson, 1972]
- Some triangulations cannot be made RI in fewer than $\Omega(n^2)$ flips



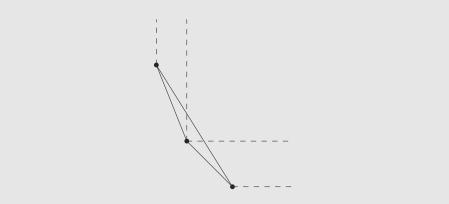
- Can we flip an arbitrary triangulation into an RI one? While getting monotonically 'closer'?
- Any triangulation can be flipped to any other in $O(n^2)$ flips [Lawson, 1972]
- Some triangulations cannot be made RI in fewer than $\Omega(n^2)$ flips



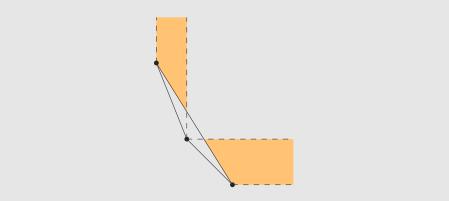
- · Count points in 'bad regions'
- No bad regions \Rightarrow the triangulation is RI

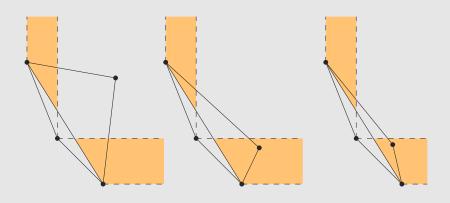


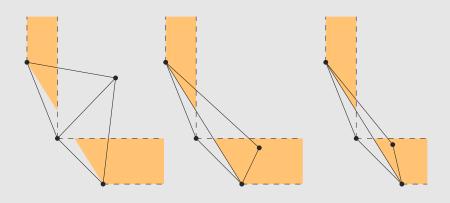
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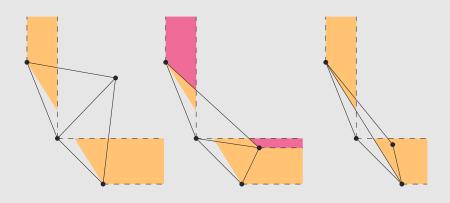


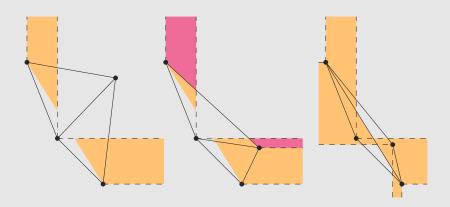
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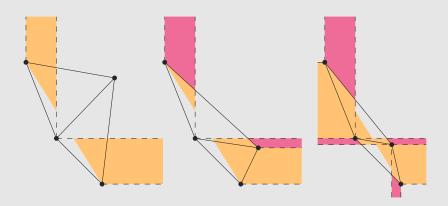






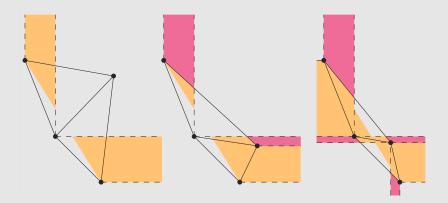






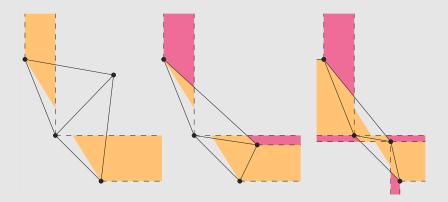
Theorem

Any augmented triangulation can be converted into an RI-triangulation with $O(n^2)$ flips.



Theorem

Any maximal triangulation can be converted into an RI-triangulation with $O(n^2)$ flips and O(n) edge deletions.



RI-Problems

- 1. RI-triangulating a polygon \checkmark
- 2. RI-triangulating a point set \checkmark
- 3. Flipping one RI-triangulation to another \checkmark
- 4. Flipping a triangulation to an RI-triangulation \checkmark

RI-Summary

- · Any polygon or point set can be RI-triangulated
- Any two RI-triangulations can be transformed into each other with $\Theta(n^2)$ flips
- Any triangulation can be transformed monotonically into an RI-triangulation with $\Theta(n^2)$ flips

