

# Proofs of Markov's Inequality & Chernoff's Bounds

**Markov:** For any non-negative r.v.  $X$ ,

$$\Pr\{X \geq t \cdot E[X]\} \leq 1/t.$$

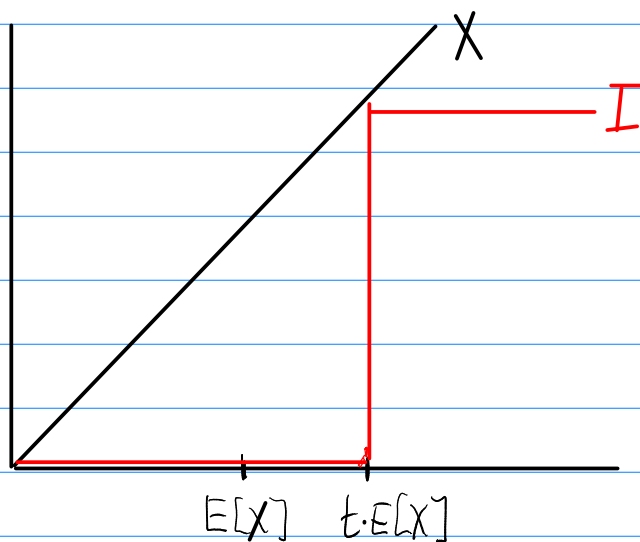
**Proof:** Let  $I = \begin{cases} t \cdot E[X] & \text{if } X \geq t \cdot E[X] \\ 0 & \text{otherwise} \end{cases}$

Notice that  $I \leq X$  . (\*)

$$E[I] = \Pr\{X \geq t \cdot E[X]\} \cdot t \cdot E[X] \leq \underbrace{E[X]}_{\text{Because of (*)}} \quad (**)$$

Divide both sides of (\*\*) by  $t \cdot E[X]$  to get

$$\Pr\{X \geq t \cdot E[X]\} \leq 1/t. \quad \text{QED}$$



Recall that events  $A$  and  $B$  are independent if and only if  $\Pr\{A \cap B\} = \Pr\{A\} \cdot \Pr\{B\}$ .

Lemma: If  $X$  and  $Y$  are independent random variables, then  
 $E[X \cdot Y] = E[X] \cdot E[Y]$ .

Proof:

$$E[X \cdot Y] = \sum_x \sum_y \Pr\{X=x \text{ and } Y=y\} \cdot xy$$

$$= \sum_x \sum_y \Pr\{X=x\} \cdot \Pr\{Y=y\} \cdot xy \quad \left[ \begin{array}{l} \text{because } X \text{ and} \\ Y \text{ are indep.} \end{array} \right]$$

$$= \sum_x \Pr\{X=x\} \cdot x \cdot \left( \sum_y \Pr\{Y=y\} \cdot y \right)$$

$$= \sum_x \Pr\{X=x\} \cdot x \cdot E[Y]$$

$$= E[Y] \cdot \sum_x \Pr\{X=x\}$$

$$= E[Y] \cdot E[X]$$

QED.

## Chernoff's Bounding Technique.

Idea: Markov's Inequality is too weak.

Instead of studying  $B$ , study  $e^{c \cdot B}$ .

If  $B$  is a little bit bigger than  $E[B]$ ,  
the  $e^{cB}$  is a lot bigger than  $E[e^{cB}]$ .

---

Let  $B = X_1 + X_2 + \dots + X_n$  where the  $X_i$ 's are independent Bernoulli( $p$ ) random variables.

$$E[e^{cX_i}] = p \cdot e^c + (1-p) \cdot e^0 = p \cdot e^c + 1 - p = 1 + p(e^c - 1)$$

$$\begin{aligned} E[e^{cB}] &= E[e^{cX_1} \cdot e^{cX_2} \cdot \dots \cdot e^{cX_n}] \\ &= E[e^{cX_1}] \cdot \dots \cdot E[e^{cX_n}] = (1 + p(e^c - 1))^n \end{aligned}$$

Next, we will use the Inequality:

$$1 + x \leq e^x \quad \text{for all } x,$$

$$\text{so } E[e^{cB}] = (1 + p(e^c - 1))^n \leq e^{p(e^c - 1)n}$$

Remember Markov:  $\Pr\{X \geq t \cdot E[X]\} \leq 1/t$ .

Same as  $\Pr\{X \geq q\} \leq \frac{E[X]}{q}$ .

If  $B \geq (1+\epsilon)E[B]$ , then  $e^{cB} \geq e^{c(1+\epsilon)np}$

$$\text{So, } \Pr\{B \geq (1+\epsilon)E[B]\} \\ = \Pr\{\exp(cB) \geq \exp(c(1+\epsilon)np)\}$$

$$\leq \frac{E[\exp(cB)]}{\exp(c(1+\epsilon)np)}$$

$$\leq \frac{\exp(p(e^c - 1)n)}{\exp(c(1+\epsilon)np)}$$

$$= \frac{1}{\exp((c(1+\epsilon) - e^c + 1)pn)}$$

$$(**) = \frac{1}{\exp((\epsilon^2 - \underbrace{e^{\frac{\epsilon^2}{1+\epsilon}}}_{\geq 0} + 1)pn)}$$

$\geq 0$  for all  $0 \leq \epsilon \leq 1$

Now we can choose  $c$ .

Let's take  $c = \epsilon^2 / (1+\epsilon)$

$$\leq \frac{1}{\exp(\epsilon^2 pn)}, \text{ for any } 0 \leq \epsilon \leq 1.$$

So, we get.

$$\Pr\{B \geq (1+\varepsilon)np\} \leq 1/e^{\varepsilon^2 np}, \text{ for any } 0 < \varepsilon < 1.$$

By choosing different values of  $c$ , we can get different bounds,

Taking  $c = \ln(1+\varepsilon)$  gives the most common version:

$$\Pr\{B \geq (1+\varepsilon)np\} \leq \left( \frac{e^\varepsilon}{(1+\varepsilon)^{1+\varepsilon}} \right)^{np}$$

Search for "Chernoff Bound" in Wikipedia for other versions.