

# Blocking Coloured Point Sets\*

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## 1 Introduction

This paper studies problems related to visibility and blocking in sets of coloured points in the plane. A point  $x$  *blocks* two points  $v$  and  $w$  if  $x$  is in the interior of the line segment  $\overline{vw}$ . Let  $P$  be a finite set of points in the plane. Two points  $v$  and  $w$  are *visible* with respect to  $P$  if no point in  $P$  blocks  $v$  and  $w$ . The *visibility graph* of  $P$  has vertex set  $P$ , where two distinct points  $v, w \in P$  are adjacent if and only if they are visible with respect to  $P$ . A point set  $B$  *blocks*  $P$  if  $P \cap B = \emptyset$  and for all distinct  $v, w \in P$  there is a point in  $B$  that blocks  $v$  and  $w$ . That is, no two points in  $P$  are visible with respect to  $P \cup B$ , or alternatively,  $P$  is an independent set in the visibility graph of  $P \cup B$ .

A set of points  $P$  is *k-blocked* if each point in  $P$  is assigned one of  $k$  colours, such that each pair of points  $v, w \in P$  are visible with respect to  $P$  if and only if  $v$  and  $w$  are coloured differently. Thus  $v$  and  $w$  are assigned the same colour if and only if some other point in  $P$  blocks  $v$  and  $w$ . We say  $P$  is  $\{n_1, \dots, n_k\}$ -blocked if it is  $k$ -blocked and for some labelling of the colours by the integers  $[k] := \{1, 2, \dots, k\}$ , the  $i$ -th colour class has exactly  $n_i$  points, for each  $i \in [k]$ . Equivalently,  $P$  is  $\{n_1, \dots, n_k\}$ -blocked if the visibility graph of  $P$  is the complete  $k$ -partite graph  $K(n_1, \dots, n_k)$ . See Figure 1 for an example.

The following fundamental conjecture regarding  $k$ -blocked point sets is the focus of this paper.

**Conjecture 1** *For each integer  $k$  there is an integer  $n$  such that every  $k$ -blocked set has at most  $n$  points.*

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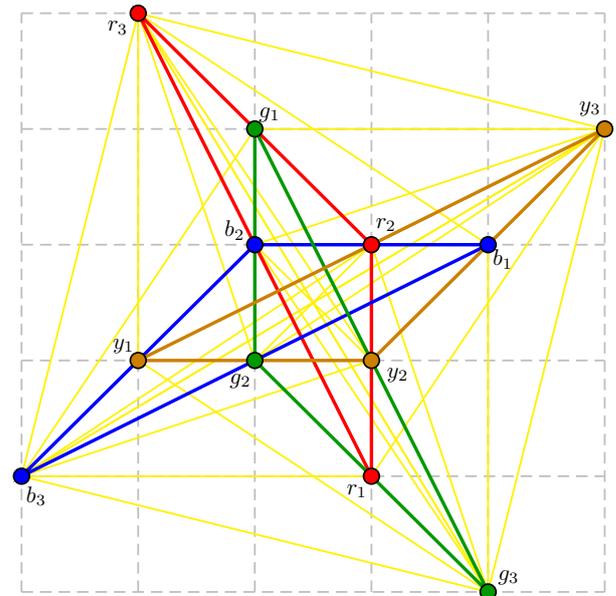


Figure 1: A  $\{3, 3, 3\}$ -blocked point set.

A  $k$ -set is a multiset of  $k$  positive integers. A  $k$ -set  $\{n_1, \dots, n_k\}$  is *representable* if there is an  $\{n_1, \dots, n_k\}$ -blocked point set. As illustrated in Figure 2, it follows from the characterisation of 2- and 3-colourable visibility graphs by Kára et al. [6] that  $\{1, 1\}$  and  $\{1, 2\}$  are the only representable 2-sets, and that  $\{1, 1, 1\}$ ,  $\{1, 1, 2\}$ ,  $\{1, 2, 2\}$  and  $\{2, 2, 2\}$  are the only representable 3-sets. In particular, every 2-blocked point set has at most 3 points, and every 3-blocked point set has at most 6 points. This proves Conjecture 1 for  $k \leq 3$ . Later we prove Conjecture 1 for  $k = 4$ .

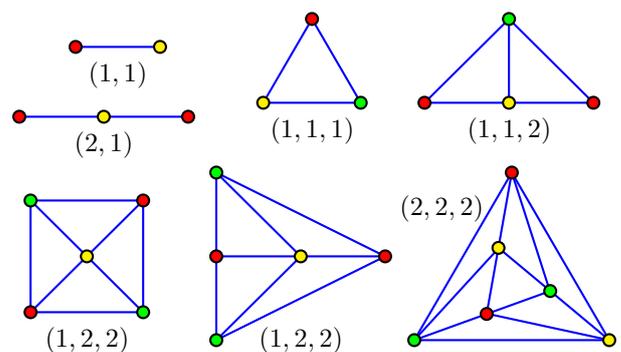


Figure 2: The 2-blocked and 3-blocked point sets.

This paper makes the following contributions. Section 2 introduces some background motivation. Section 3 describes methods for constructing  $k$ -blocked sets from a given  $(k - 1)$ -blocked set. These methods lead to a characterisation of representable  $k$ -sets when each colour class has at most three points. Section 4 studies the  $k = 4$  case in more detail. In particular, we characterise the representable 4-sets, and conclude that the example in Figure 1 is in fact the largest 4-blocked point set. Section 5 introduces a special class of  $k$ -blocked sets (so-called midpoint-blocked sets) that lead to a construction of the largest known  $k$ -blocked sets for infinitely many values of  $k$ .

Also note the following easily proved properties.

**Lemma 1 ([2])** *At most three points are collinear in every  $k$ -blocked point set.*

**Lemma 2 ([2])** *Each colour class in a  $k$ -blocked point set is in general position (no three collinear).*

## 2 Some Background Motivation

Much recent research on blockers began with the following conjecture by Kára et al. [6].

**Conjecture 2 (Big-Line-Big-Clique Conjecture [6])** *For all integers  $t$  and  $\ell$  there is an integer  $n$  such that for every finite set  $P$  of at least  $n$  points in the plane:*

- $P$  contains  $\ell$  collinear points, or
- $P$  contains  $t$  pairwise visible points (that is, the visibility graph of  $P$  contains a  $t$ -clique).

Conjecture 2 is true for  $t \leq 5$ , but is open for  $t \geq 6$  or  $\ell \geq 4$ ; see [10, 1]. Jan Kára suggested the following weakening of Conjecture 2.

**Conjecture 3 ([10])** *For all integers  $t$  and  $\ell$  there is an integer  $n$  such that for every finite set  $P$  of at least  $n$  points in the plane:*

- $P$  contains  $\ell$  collinear points, or
- the chromatic number of the visibility graph of  $P$  is at least  $t$ .

Clearly Conjecture 2 implies Conjecture 3.

**Proposition 3** *Conjecture 3 with  $\ell = 4$  and  $t = k + 1$  implies Conjecture 1.*

**Proof.** Assume Conjecture 3 holds for  $\ell = 4$  and  $t = k + 1$ . Suppose  $P$  is a  $k$ -blocked set of at least  $n$  points. By Lemma 1, at most three points are collinear. Thus the first conclusion of Conjecture 3 does not hold. Since the visibility graph of  $P$  is  $k$ -colourable, the second conclusion of Conjecture 3 does not hold. This contradiction proves that every  $k$ -blocked set has less than  $n$  points, and Conjecture 1 holds.  $\square$

Since Conjecture 2 holds for  $t \leq 5$ , Conjecture 1 holds for  $k \leq 4$ . Let  $b(n)$  be the minimum integer such that some set of  $n$  points in the plane in general position is blocked by some set of  $b(n)$  points. Linear lower bounds on  $b(n)$  are known [7, 3], but many authors have conjectured or stated as an open problem that  $b(n)$  is super-linear.

**Conjecture 4 ([7, 9, 3, 10])**  $\frac{b(n)}{n} \rightarrow \infty$  as  $n \rightarrow \infty$ .

Pór and Wood [10] proved that Conjecture 4 implies Conjecture 3, and thus implies Conjecture 1. That Conjecture 1 is implied by a number of other well-known conjectures, yet remains challenging, adds to its interest.

## 3 $k$ -Blocked Sets with Small Colour Classes

We now describe some methods for building blocked point sets from smaller blocked point sets.

**Lemma 4** *Let  $G$  be a visibility graph. Let  $i \in \{1, 2, 3\}$ . Furthermore suppose that if  $i \geq 2$  then  $V(G) \neq \emptyset$ , and if  $i = 3$  then not all the vertices of  $G$  are collinear. Let  $G_i$  be the graph obtained from  $G$  by adding an independent set of  $i$  new vertices, each adjacent to every vertex in  $G$ . Then  $G_1$ ,  $G_2$ , and  $G_3$  are visibility graphs.*

**Proof.** For distinct points  $p$  and  $q$ , let  $\overleftarrow{pq}$  denote the ray that is (1) contained in the line through  $p$  and  $q$ , (2) starting at  $p$ , and (3) not containing  $q$ . Let  $\mathcal{L}$  be the union of the set of lines containing at least two vertices in  $G$ .

$i = 1$ : Since  $\mathcal{L}$  is the union of finitely many lines, there is a point  $p \notin \mathcal{L}$ . Thus  $p$  is visible from every vertex of  $G$ . By adding a new vertex at  $p$ , we obtain a representation of  $G_1$  as a visibility graph.

$i = 2$ : Let  $p$  be a point not in  $\mathcal{L}$ . Let  $v$  be a vertex of  $G$ . Each line in  $\mathcal{L}$  intersects  $\overleftarrow{vp}$  in at most one point. Thus  $\overleftarrow{vp} \setminus \mathcal{L} \neq \emptyset$ . Let  $q$  be a point in  $\overleftarrow{vp} \setminus \mathcal{L}$ . Thus  $p$  and  $q$  are visible from every vertex of  $G$ , but  $p$  and  $q$  are blocked by  $v$ . By adding new vertices at  $p$  and  $q$ , we obtain a representation of  $G_2$  as a visibility graph.

$i = 3$ : Let  $u, v, w$  be non-collinear vertices in  $G$ . Let  $p$  be a point not in  $\mathcal{L}$  and not in the convex hull of  $\{u, v, w\}$ . Without loss of generality,  $\overline{uv} \cap \overline{vw} \neq \emptyset$ . There are infinitely many pairs of points  $q \in \overleftarrow{up}$  and  $r \in \overleftarrow{vp}$  such that  $w$  blocks  $q$  and  $r$ . Thus there are such  $q$  and  $r$  both not in  $\mathcal{L}$ . By construction,  $u$  blocks  $p$  and  $q$ , and  $v$  blocks  $p$  and  $r$ . By adding new vertices at  $p$ ,  $q$  and  $r$ , we obtain a representation of  $G_3$  as a visibility graph.  $\square$

We now characterise the representable ( $\geq 4$ )-sets, assuming each colour class has at most three points.

**Proposition 5** A  $k$ -set  $\{n_1, \dots, n_k\}$  is representable whenever  $k \geq 4$  and each  $n_i \leq 3$ , except for  $\{1, 3, 3, 3\}$  which is not representable [2].

**Proof.** We say  $\{n_1, \dots, n_k\}$  contains  $\{n_1, \dots, n_{i-1}, n_{i+1}, \dots, n_k\}$  for each  $i \in [k]$ . We proceed by induction on  $k$ . If  $\{n_1, \dots, n_k\}$  contains a representable  $(k-1)$ -set, then  $\{n_1, \dots, n_k\}$  is also representable by Lemma 4. (Since  $k \geq 4$  the assumptions in Lemma 4 hold.) Now assume that every  $(k-1)$ -set contained in  $\{n_1, \dots, n_k\}$  is not representable. By induction, we may assume that  $k \leq 5$ . Moreover, if  $k = 5$  then  $\{n_1, \dots, n_5\}$  must contain  $\{1, 3, 3, 3\}$  (since by induction all other 4-sets are representable). Similarly, if  $k = 4$  then  $\{n_1, \dots, n_4\}$  must contain  $\{1, 1, 3\}$ ,  $\{1, 2, 3\}$ ,  $\{1, 3, 3\}$ ,  $\{2, 2, 3\}$ ,  $\{2, 3, 3\}$  or  $\{3, 3, 3\}$  (since  $\{1, 1, 1\}$ ,  $\{1, 1, 2\}$ ,  $\{1, 2, 2\}$  and  $\{2, 2, 2\}$  are representable). The following table describes the construction in each case.

$\{1, 1, 1, x\}$	contains $\{1, 1, 1\}$
$\{1, 1, 2, x\}$	contains $\{1, 1, 2\}$
$\{1, 1, 3, 3\}$	Figure 1 minus $\{r_1, g_3, r_3, g_1\}$
$\{1, 2, 2, x\}$	contains $\{1, 2, 2\}$
$\{1, 2, 3, 3\}$	Figure 1 minus $\{g_1, g_3, r_3\}$
$\{2, 2, 2, x\}$	contains $\{2, 2, 2\}$
$\{2, 2, 3, 3\}$	Figure 1 minus $\{g_3, r_3\}$
$\{2, 3, 3, 3\}$	Figure 1 minus $g_3$
$\{1, 1, 3, 3, 3\}$	contains $\{1, 1, 3, 3\}$
$\{1, 2, 3, 3, 3\}$	contains $\{1, 2, 3, 3\}$
$\{1, 3, 3, 3, 3\}$	contains $\{3, 3, 3, 3\}$

□

#### 4 4-Blocked Point Sets

As shown in Section 2, Conjecture 1 holds for  $k \leq 4$ . That is, every 4-blocked set has bounded size. An explicit bound of  $2^{790}$  follows from a result of Abel et al. [1], which can be improved to  $2^{578}$  using a recent result by Dumitrescu et al. [3]; see [2]. Before characterising all representable 4-sets, we give a simple proof that every 4-blocked point set has bounded size.

**Proposition 6** Every 4-blocked set has at most 36 points.

**Proof.** Let  $P$  be a 4-blocked set. Suppose that  $|P| \geq 37$ . Let  $S$  be the largest colour class. Thus  $|S| \geq 10$ . By Lemma 2,  $S$  is in general position. By a theorem of Harborth [4], some 5-point subset  $K \subseteq S$  is the vertex-set of an empty convex pentagon  $\text{conv}(K)$ . Let  $T := P \cap (\text{conv}(K) - K)$ . Since  $\text{conv}(K)$  is empty with respect to  $S$ , each point in  $T$  is not in  $S$ . Thus  $T$  is 3-blocked.  $K$  needs at least 8 blockers (5 blockers for the edges on the boundary of  $\text{conv}(K)$ , and 3 blockers for the chords of  $\text{conv}(K)$ ). Thus  $|T| \geq 8$ . But every 3-blocked set has at most 6 points, which is a contradiction. Hence  $|P| \leq 36$ . □

**Theorem 7** A 4-set  $\{a, b, c, d\}$  is representable if and only if:

- $\{a, b, c, d\} = (4, 2, 2, 1)$ , or
- $\{a, b, c, d\} = (4, 2, 2, 2)$ , or
- all of  $a, b, c, d \leq 3$  except for  $\{3, 3, 3, 1\}$

**Proof Sketch.** Figure 3 shows  $\{4, 2, 2, 1\}$ -blocked and  $\{4, 2, 2, 2\}$ -blocked point sets. When  $a, b, c, d \leq 3$ , the required constructions are described in Proposition 5. Now we prove the converse. Let  $P$  be a 4-blocked point set. We prove [2] that if some colour class  $S$  contains a 4-point subset  $K$ , such that  $\text{conv}(K)$  is a convex quadrilateral that is empty with respect to  $S$ , then  $P$  is  $\{4, 2, 2, 1\}$ -blocked. Moreover, if some colour class  $S$  has at least five points, then by Lemma 2 and a theorem of Esther Klein,  $S$  contains such a subset  $K$ —implying  $P$  is  $\{4, 2, 2, 1\}$ -blocked, which is a contradiction. Hence each colour class has at most four points. Let  $S$  be a largest colour class. If  $S$  consists of four points in convex position, then  $P$  is  $\{4, 2, 2, 1\}$ -blocked (just set  $K := S$ ). If  $S$  consists of four points in nonconvex position, then we prove [2] that  $P$  is  $\{4, 2, 2, 2\}$ -blocked. Otherwise  $|S| \leq 3$ , and we are done by Proposition 5. □

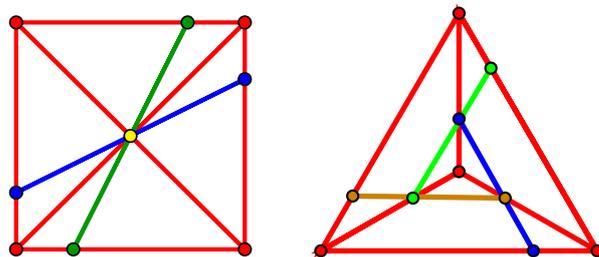


Figure 3:  $\{4, 2, 2, 1\}$ -blocked and  $\{4, 2, 2, 2\}$ -blocked point sets.

**Corollary 8** Every 4-blocked set has at most 12 points, and there is a 4-blocked set with 12 points.

#### 5 Midpoint-Blocked Point Sets

A  $k$ -blocked point set  $P$  is  $k$ -midpoint-blocked if for each monochromatic pair of distinct points  $v, w \in P$  the midpoint of  $\overline{vw}$  is in  $P$ . Of course, the midpoint of  $\overline{vw}$  blocks  $v$  and  $w$ . A point set  $P$  is  $\{n_1, \dots, n_k\}$ -midpoint-blocked if it is  $\{n_1, \dots, n_k\}$ -blocked and  $k$ -midpoint-blocked. For example, the point set in Figure 1 is  $\{3, 3, 3, 3\}$ -midpoint-blocked.

Another interesting example is the projection<sup>1</sup> of  $[3]^d$ . With  $d = 1$  this point set is  $\{2, 1\}$ -blocked, with  $d = 2$  it is  $\{4, 2, 2, 1\}$ -blocked, and with  $d = 3$  it is  $\{8, 4, 4, 4, 2, 2, 2, 1\}$ -blocked. In general, each set of

<sup>1</sup>If  $G$  is the visibility graph of some point set  $P \subseteq \mathbb{R}^d$ , then  $G$  is the visibility graph of some projection of  $P$  to  $\mathbb{R}^2$ .

points with exactly the same set of coordinates equal to 2 is a colour class, there are  $2^{d-i}$  colour classes of points with exactly  $i$  coordinates equal to 2, and  $[3]^d$  is  $\{\binom{d}{i} \times 2^i : i \in [0, d]\}$ -midpoint-blocked and  $2^d$ -midpoint-blocked.

Hernández-Barrera et al. [5] defined  $m(n)$  to be the minimum number of midpoints determined by some set of  $n$  points in general position in the plane, and proved that  $m(n) \leq cn^{\log_2 3} = cn^{1.585\dots}$ . This upper bound was improved by Pach [8] (and later by Matousek [7]) to  $m(n) \leq nc^{\sqrt{\log n}}$ . Hernández-Barrera et al. [5] conjectured that  $m(n)$  is super-linear, which was verified by Pach [8]; that is,  $\frac{m(n)}{n} \rightarrow \infty$  as  $n \rightarrow \infty$ . Pór and Wood [10] proved the following more precise version: For some  $c > 0$ , for all  $\epsilon > 0$  there is an integer  $N(\epsilon)$  such that  $m(n) \geq cn(\log n)^{1/(3+\epsilon)}$  for all  $n \geq N(\epsilon)$ .

**Theorem 9** For each  $k$  there is an integer  $n$  such that every  $k$ -midpoint-blocked set has at most  $n$  points. More precisely, there is an absolute constant  $c$  and for each  $\epsilon > 0$  there is an integer  $N(\epsilon)$ , such that for all  $k$ , every  $k$ -midpoint-blocked set has at most  $k \max\{N(\epsilon), c^{(k-1)^{3+\epsilon}}\}$  points.

**Proof.** Let  $P$  be  $k$ -midpoint-blocked set of  $n$  points. We may assume that  $\frac{n}{k} > N(\epsilon)$ . Let  $S$  be a set of exactly  $s := \lceil \frac{n}{k} \rceil$  monochromatic points in  $P$ . Thus  $S$  is in general position by Lemma 2. And for every pair of distinct points  $v, w \in S$  the midpoint of  $\overline{vw}$  is in  $P - S$ . Thus  $c \frac{n}{k} (\log \frac{n}{k})^{1/(3+\epsilon)} \leq m(s) \leq n - s \leq n(1 - \frac{1}{k})$ . Hence  $(\log \frac{n}{k})^{1/(3+\epsilon)} \leq (k-1)/c$ , implying  $n \leq k2^{((k-1)/c)^{3+\epsilon}}$ . The result follows.  $\square$

We now construct  $k$ -midpoint-blocked point sets with a ‘large’ number of points. The method is based on the following product of point sets  $P$  and  $Q$ . Let  $(x_v, y_v)$  be the coordinates of each  $v \in P \cup Q$ . Let  $P \times Q$  be the point set  $\{(v, w) : v \in P, w \in Q\}$  where  $(v, w)$  is at  $(x_v, y_v, x_w, y_w)$  in 4-dimensional space. For brevity we do not distinguish between a point in  $\mathbb{R}^4$  and its image in an occlusion-free projection of the visibility graph of  $P \times Q$  into  $\mathbb{R}^2$ .

**Lemma 10** If  $P$  is a  $\{n_1, \dots, n_k\}$ -midpoint-blocked point set and  $Q$  is a  $\{m_1, \dots, m_\ell\}$ -midpoint-blocked point set, then  $P \times Q$  is  $\{n_i m_j : i \in [k], j \in [\ell]\}$ -midpoint-blocked.

**Proof.** Colour each point  $(v, w)$  in  $P \times Q$  by the pair  $(\text{col}(v), \text{col}(w))$ . There are  $n_i m_j$  points for the  $(i, j)$ -th colour class. It is straightforward to verify that two points in  $P \times Q$  are blocked if and only if they have the same colour. Thus  $P \times Q$  is blocked. Since every blocker is a midpoint,  $P \times Q$  is midpoint-blocked.  $\square$

Say  $P$  is a  $k$ -midpoint blocked set of  $n$  points. By Lemma 10, the  $i$ -fold product  $P^i := P \times \dots \times P$  is a  $k^i$ -blocked set of  $n^i = (k^i)^{\log_k n}$  points. Taking  $P$  to be

the  $\{3, 3, 3, 3\}$ -midpoint-blocked point set in Figure 1, we obtain the following result, which describes the largest known construction of  $k$ -blocked point sets.

**Theorem 11** For all  $k$  a power of 4, there is a  $k$ -blocked set of  $k^{\log_4 12} = k^{1.79\dots}$  points.

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